

Tolerating Slowdowns in Replicated State Machines using Copilots

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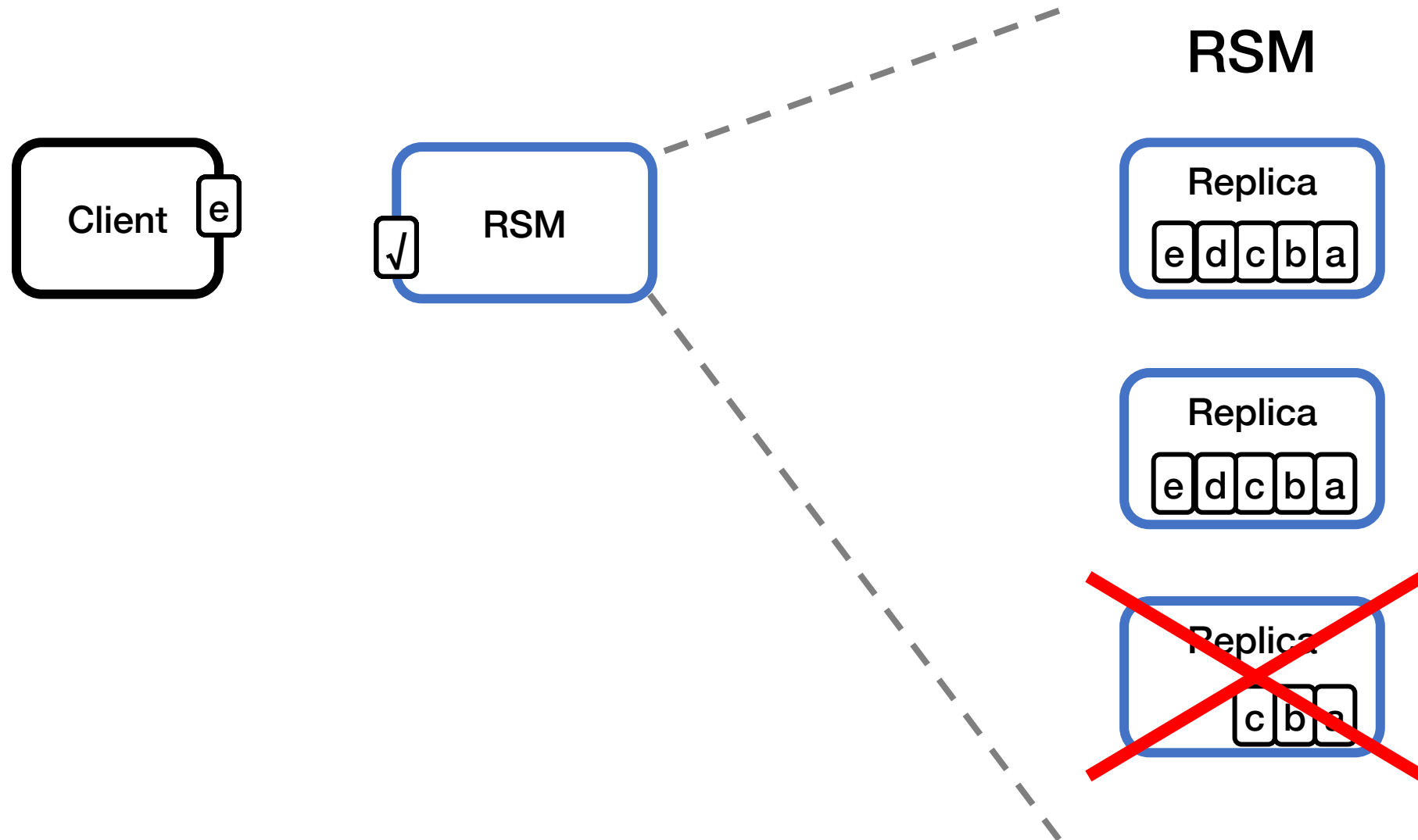
Replicated State Machine (RSM)

- Fault tolerant group of replicas that acts like a single machine that does not fail
- RSMs are everywhere!
 - Distributed database, cloud storage, coordination services, ...



Cloud
Spanner

Fault Tolerance for High Availability

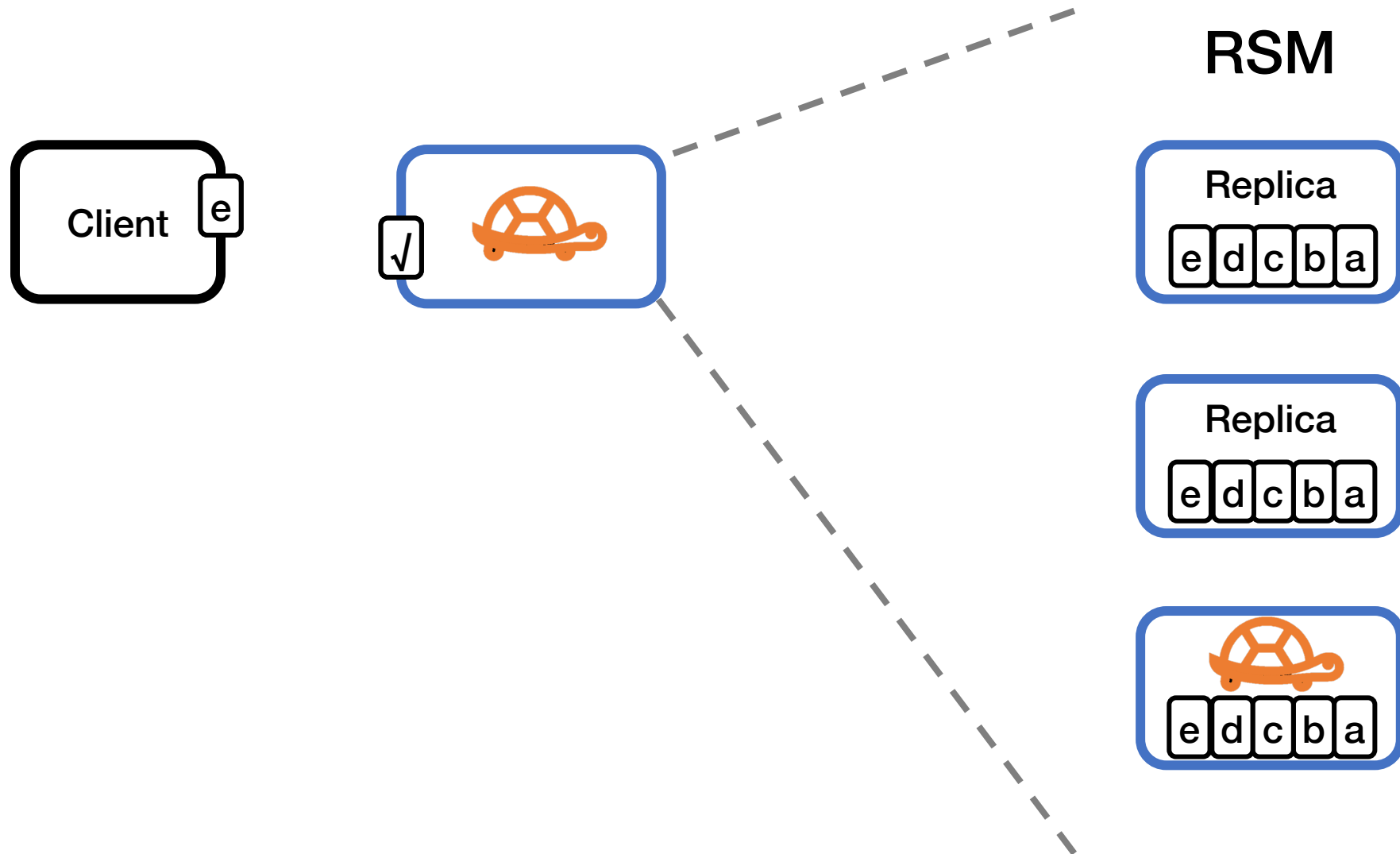


Replicas can Slowdown instead of Fail

- **Many causes:**
 - Misconfigurations
 - Partial hardware failures
 - Garbage collection events
 - ...
- **Effect: Replica takes longer than usual to send responses**

RSMs tolerate failures, not slowdowns

Slowdowns Hurt Availability



We need slowdown tolerance!

Slowdowns Take Different Forms

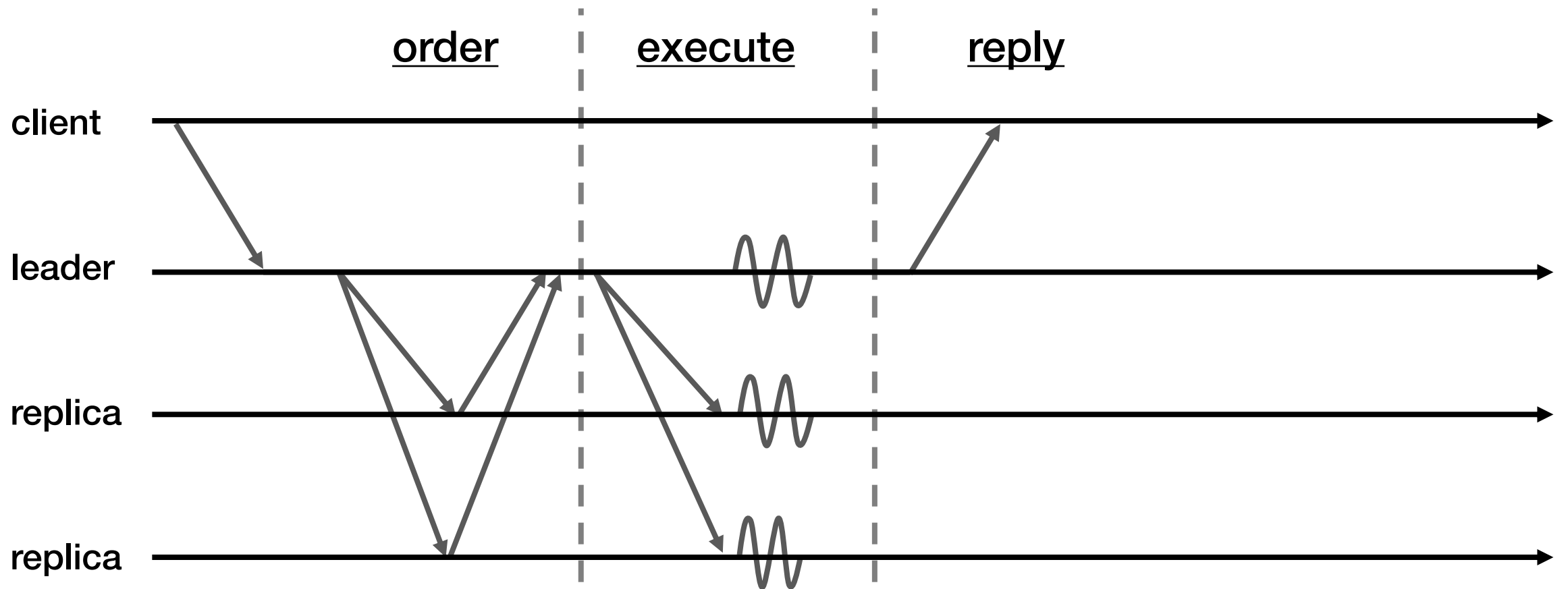
- **Duration**
 - Transient slowdowns: not handled in general
 - Long-term slowdowns: eventually detected, but need to tolerate between onset and end of reconfiguration
- **Severity**
 - 10ms additional delay or 80ms?
- **Scope**
 - All processing paths or a subset?

Defining Slowdown Tolerance

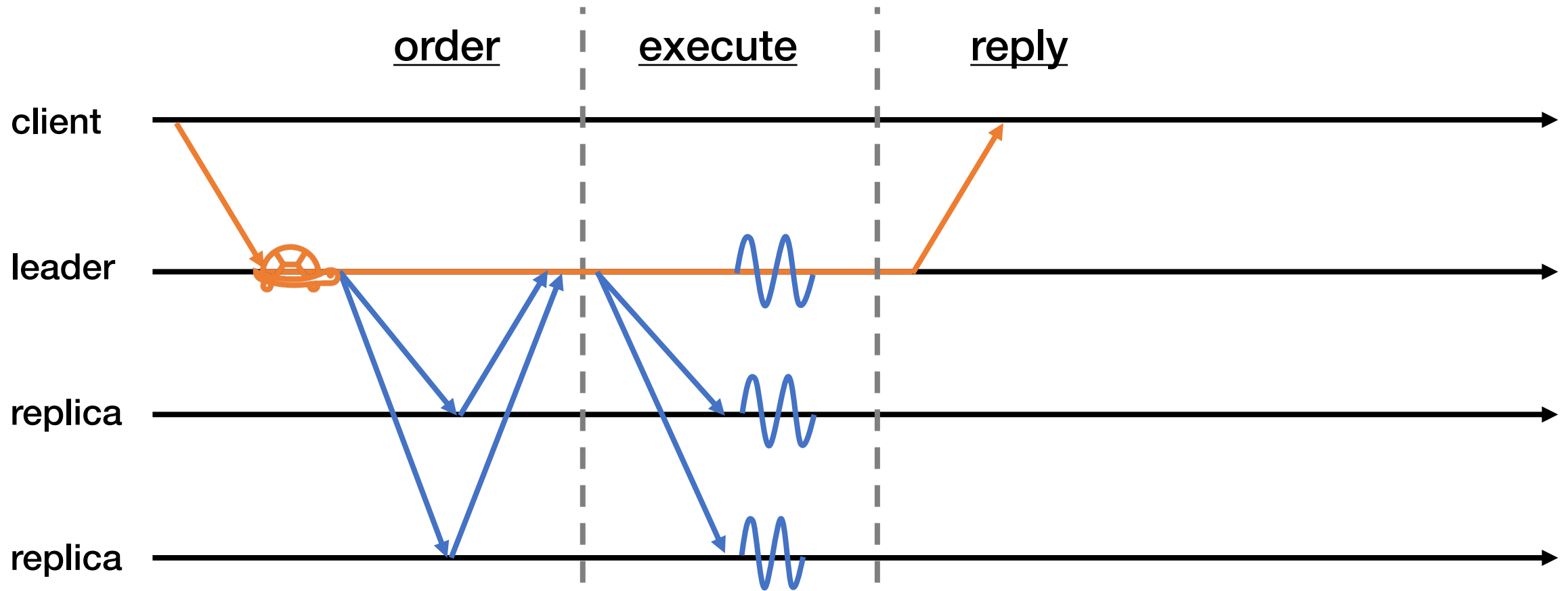
- “slow” replica = responses to messages take more than threshold time t over normal response time
- An RSM is s -slowdown-tolerant if it is not slow despite s slow replicas
 - Replacing the s slowest replicas with normal replicas should not change performance much
- This work’s focus: 1-slowdown-tolerance

**No existing consensus protocol
is 1-slowdown-tolerant**

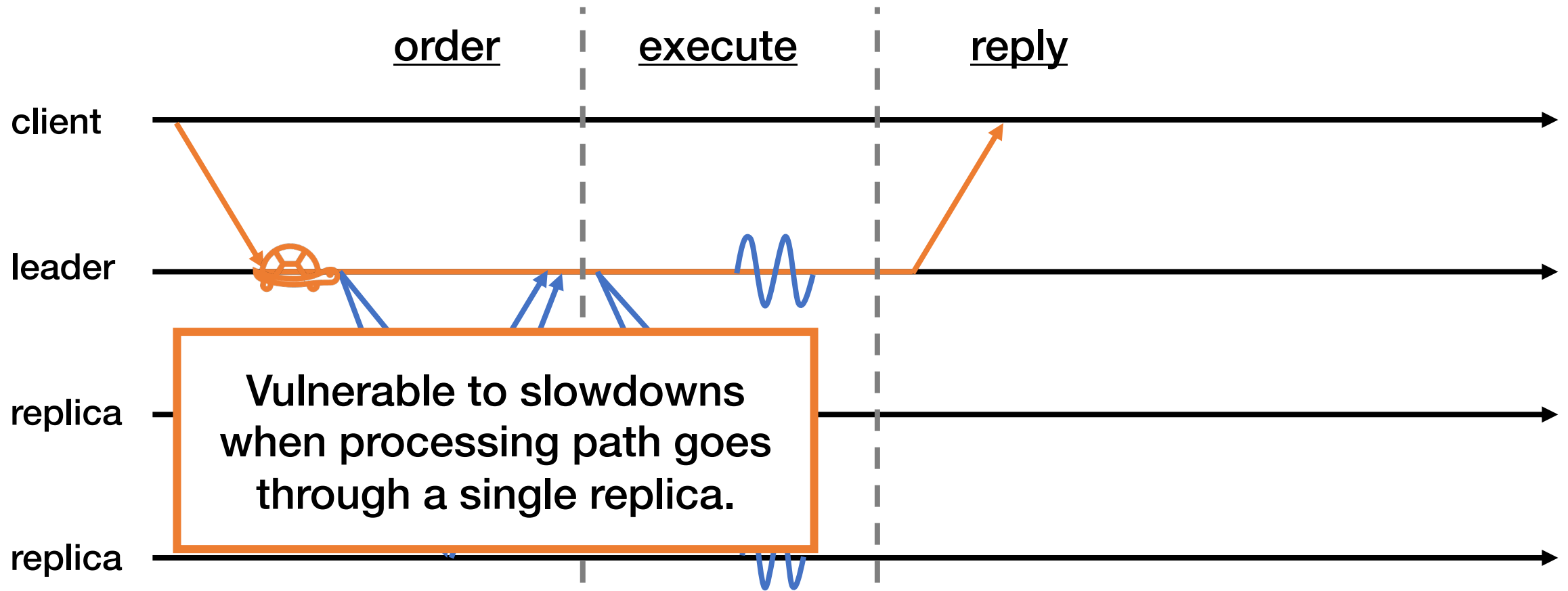
Multi-Paxos is Not 1-Slowdown-Tolerant



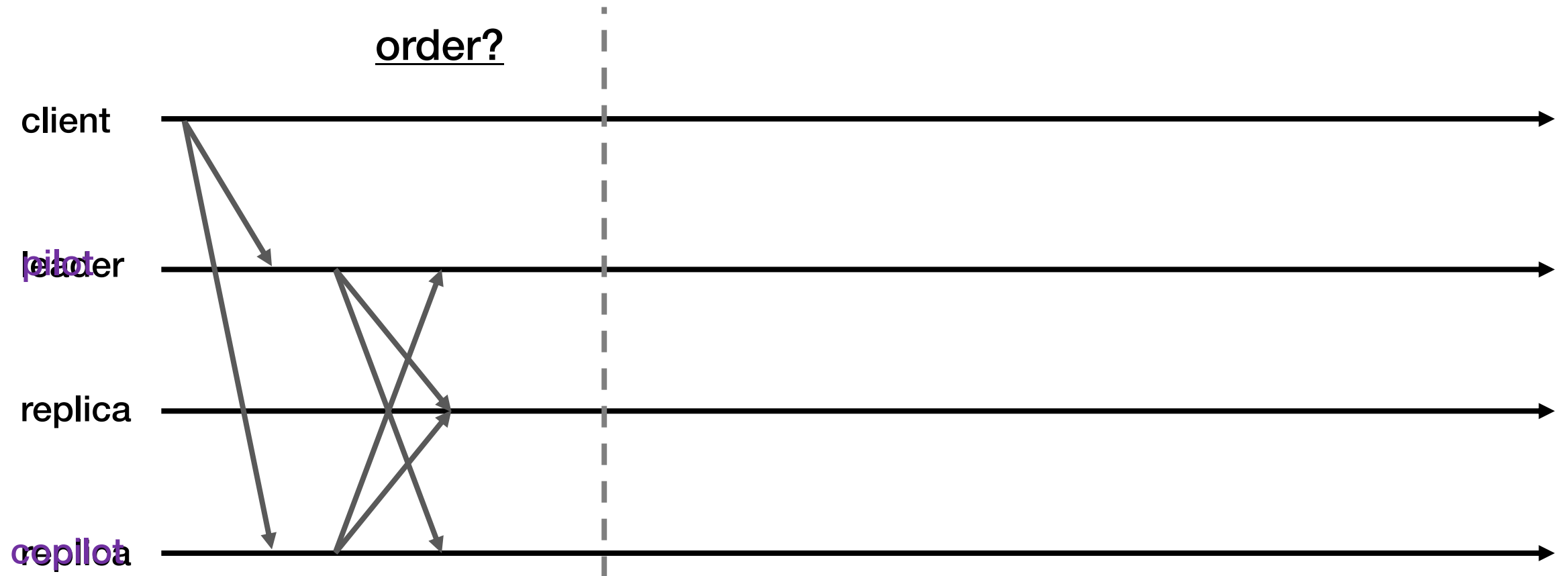
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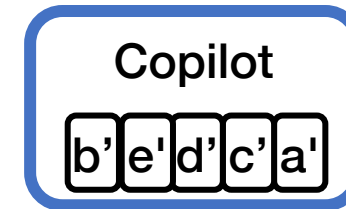
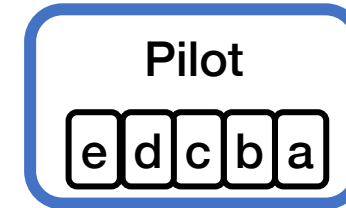


Copilot: First 1-Slowdown-Tolerant Protocol



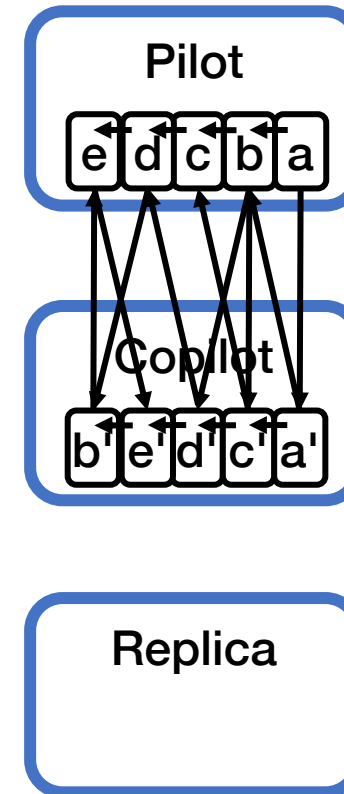
Ordering: Use Two Logs

RSM



Ordering: Combine Logs with Dependencies

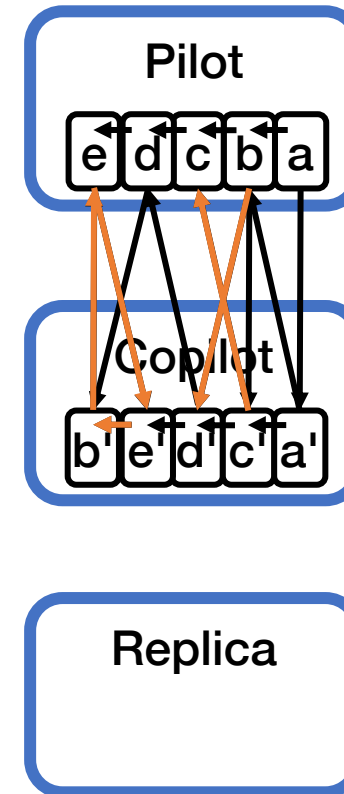
RSM



Ordering: Dependency Cycles

Break cycles deterministically

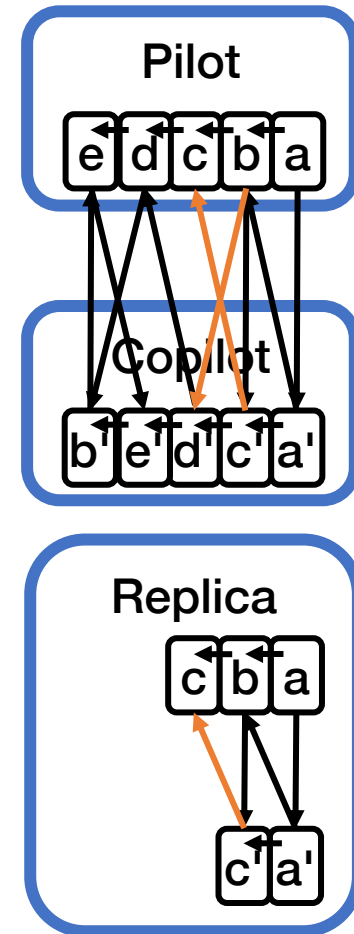
RSM



Ordering: A Tricky Case

Possible ordering:
a, a', b, c', c

RSM

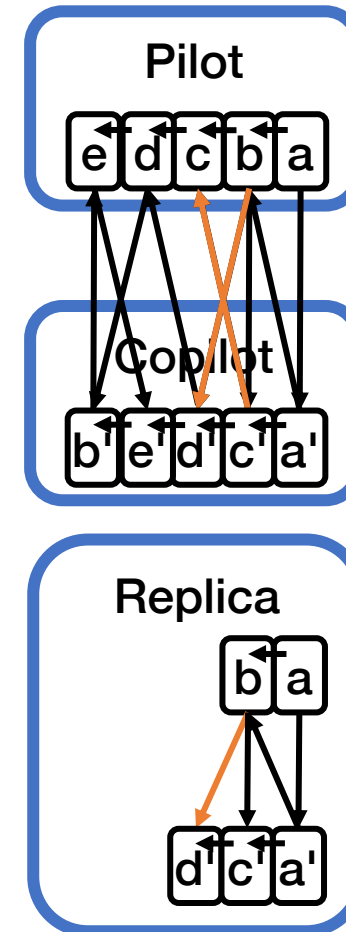


Ordering: A Tricky Case

~~Possible ordering:
a, a', b, c', c~~

~~Possible ordering:
a, a', b, c', d'~~

RSM

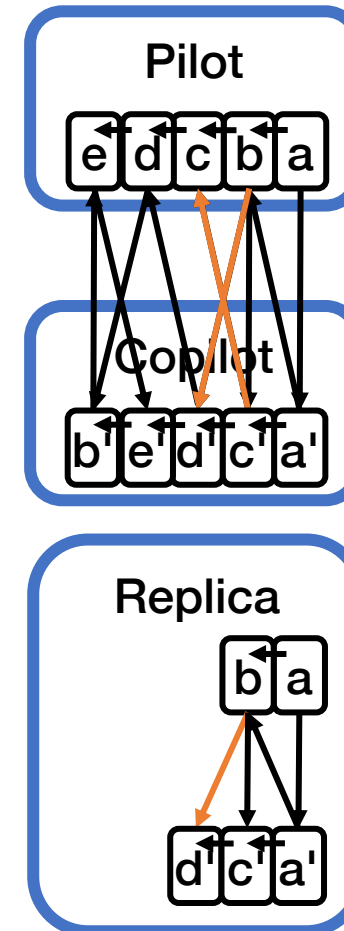


Ordering: Same on All Replicas

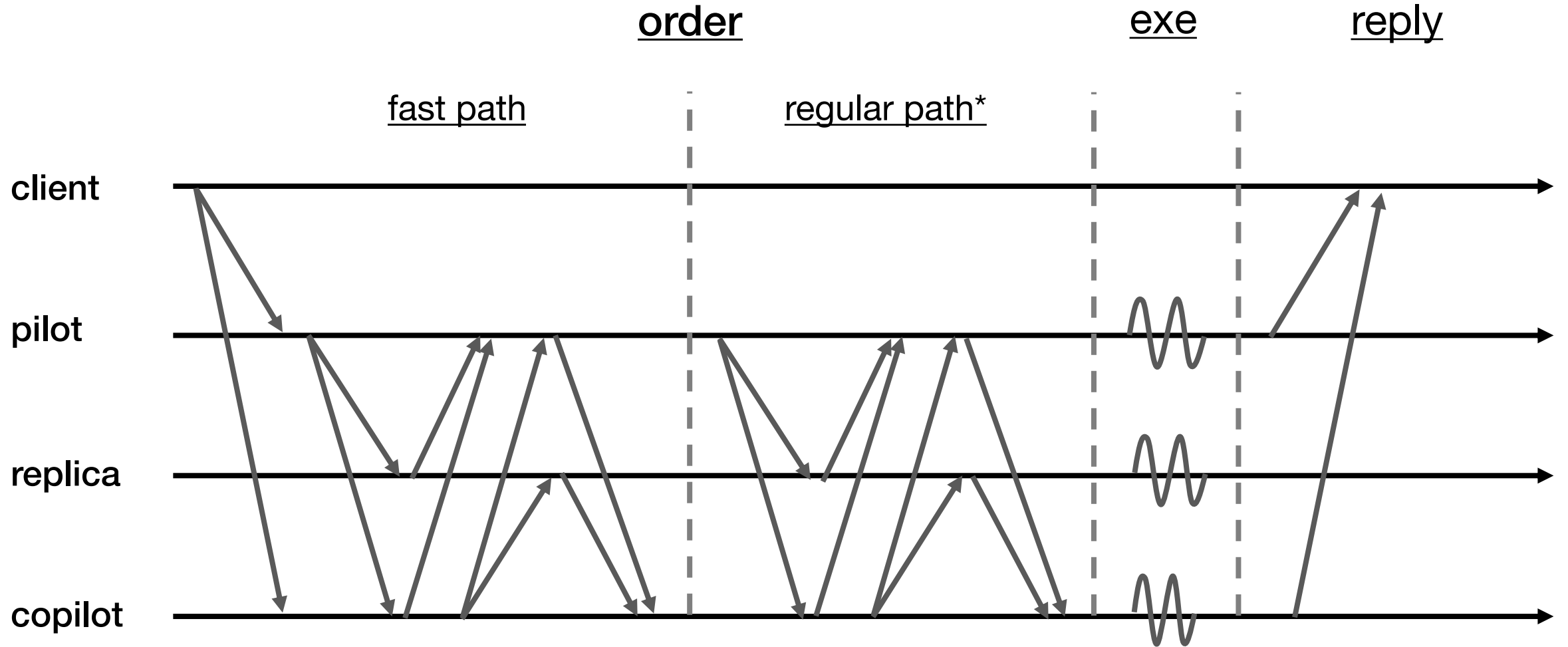
Compatibility check:
Only accept dependency if it
cannot lead to multiple orders

Break cycles deterministically

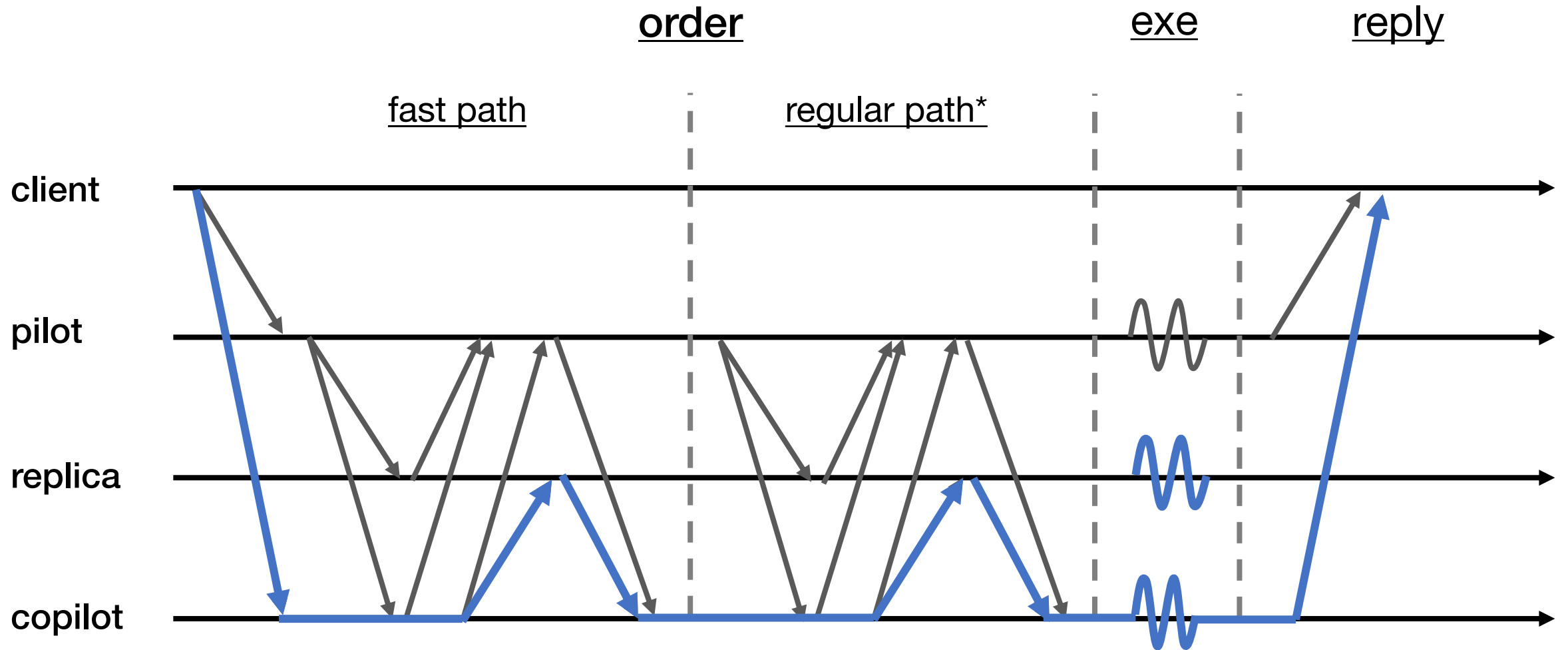
RSM



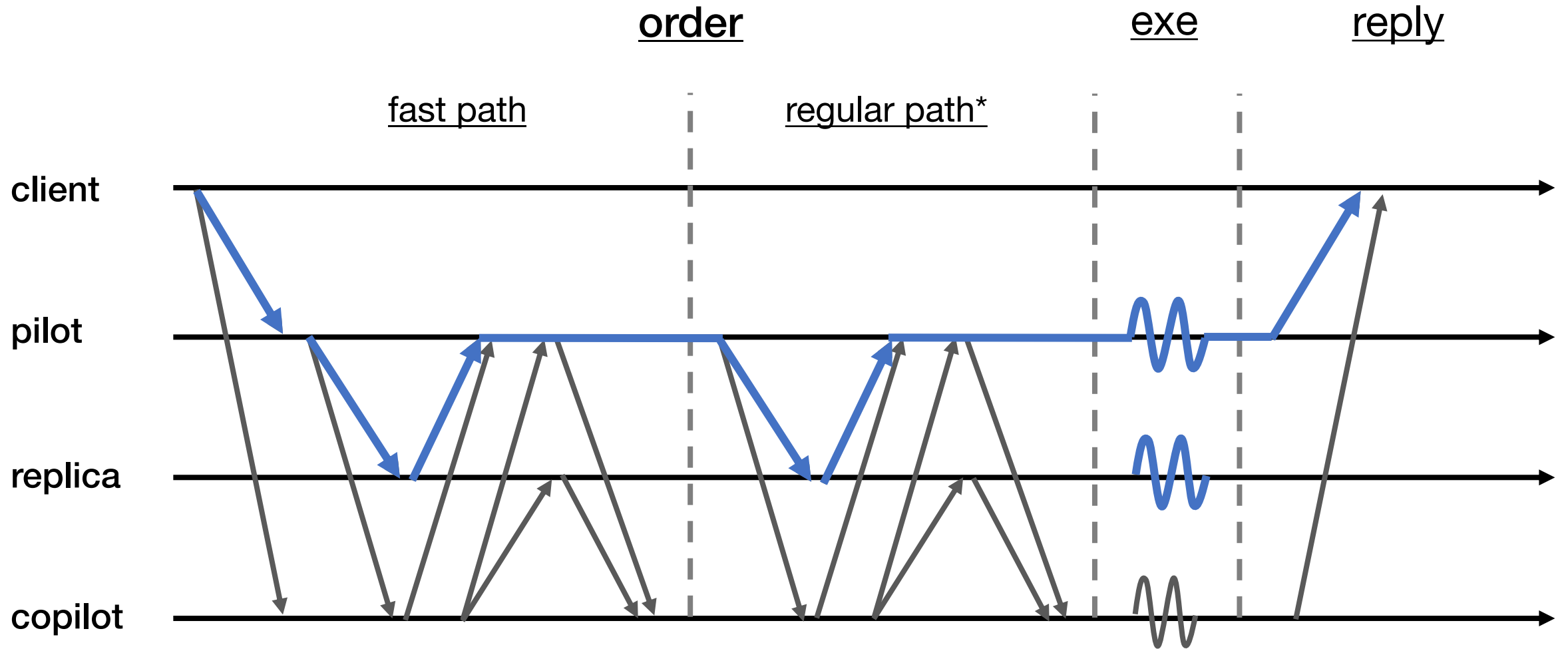
Copilot Protocol



Copilot Protocol

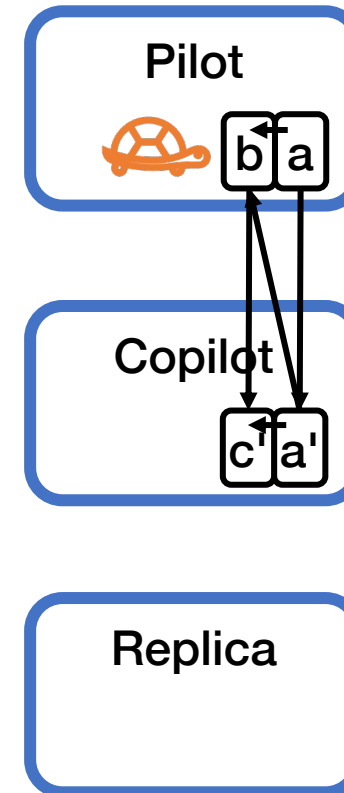


Copilot Protocol



Copilot Protocol: Dependencies?

RSM



Solution: fast takeover the slow pilot's ordering work!

Copilot Protocol: Summary

- **Proactive redundancy: two pilots process all commands**
- **Use dependencies to combine ordering from two pilots**
 - **Compatibility check ensures same order on all replicas**
 - **Cycles broken by priority**
 - **Fast takeover to avoiding waiting on slow pilot**

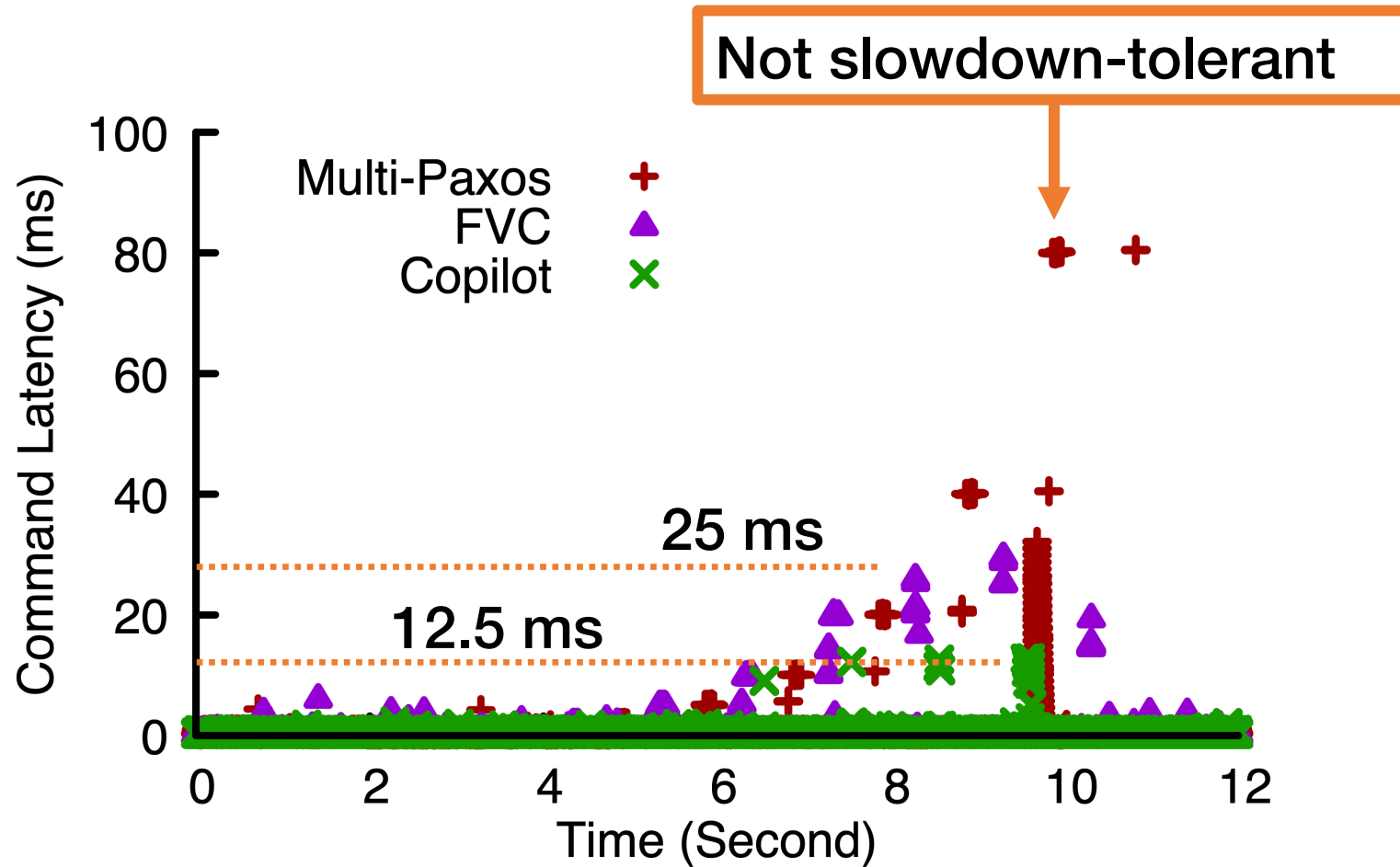
Optimizations

- **Ping-Pong Batching**
 - Improve Copilot's performance when both pilots are fast
 - Pilots propose compatible orderings and commit on fast path
- **Null Dependency Elimination**
 - Improve Copilot's performance when one pilot is slow
 - Allow a fast pilot to safely avoid waiting on commits from a continually slow pilot and thus avoid fast takeover

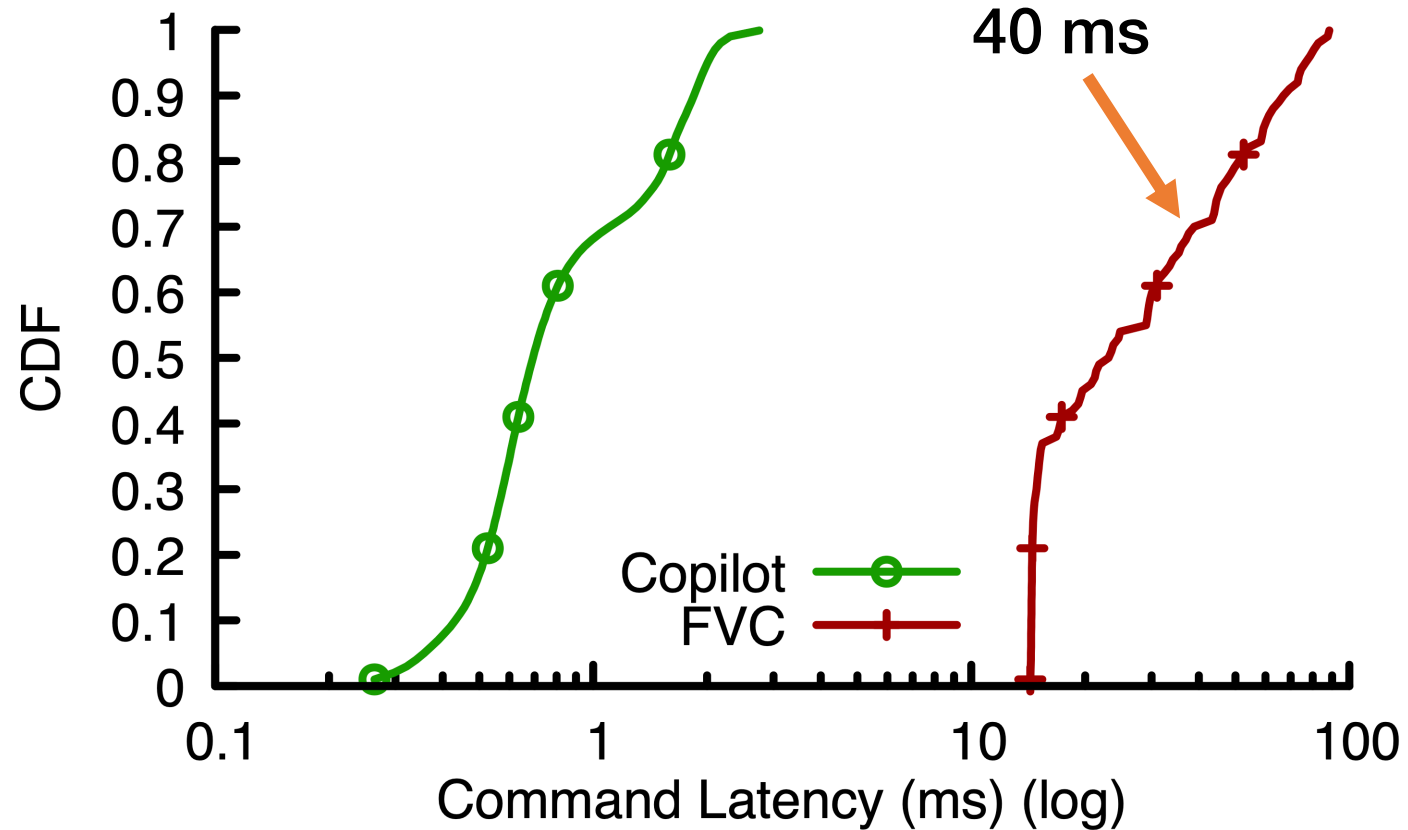
Evaluation

- Tolerate slowdowns that are transient, have varying manifestations, have varying severity?
- How does Copilot perform without slow replicas?
- 5-replica RSM, moderate load
- Replicas and clients in the same datacenter
- Baselines:
 - EPaxos
 - Multi-Paxos
 - Fast-View-Change (10 ms view-change timeout)

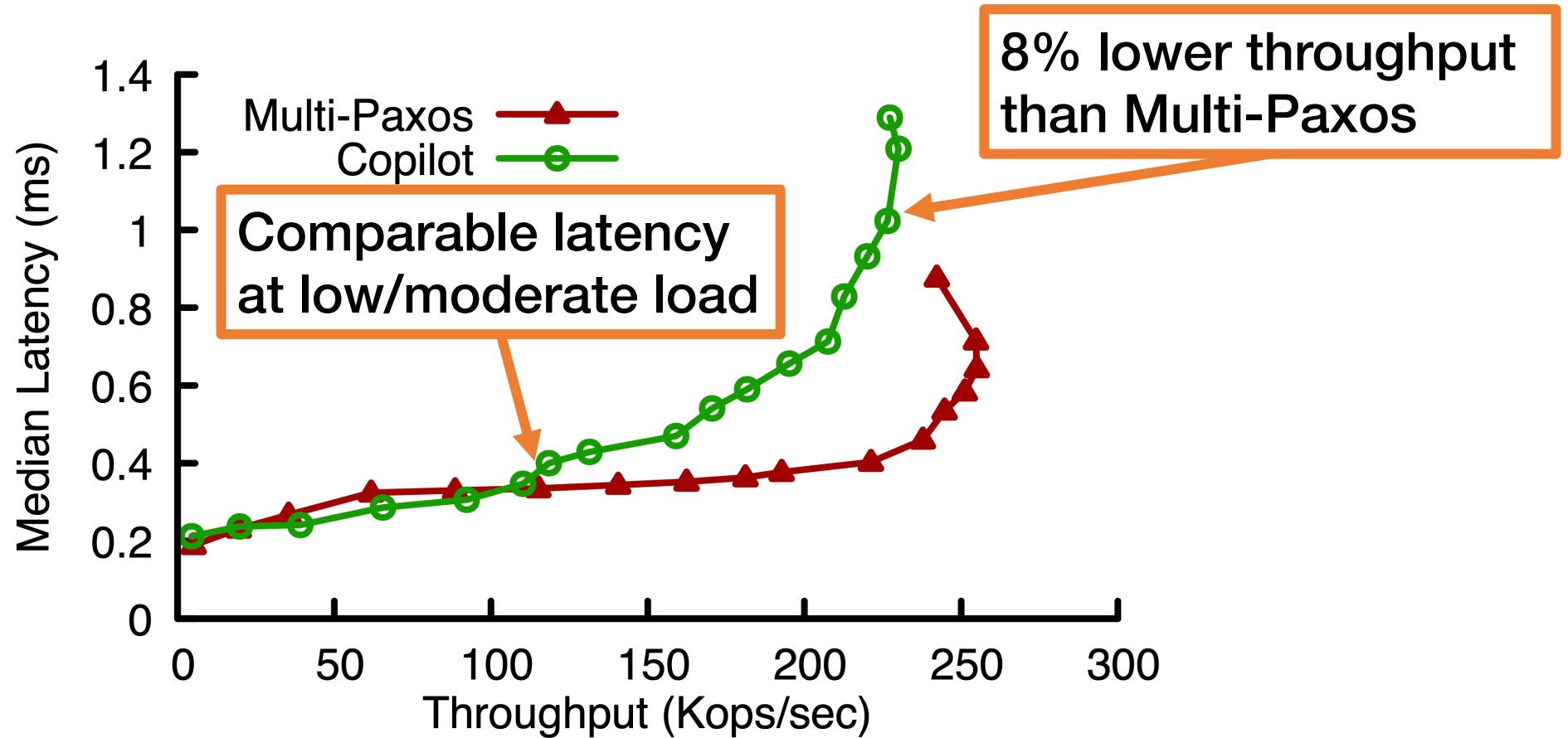
Transient Slowdowns



Gradual Slowdown



Performance Without Slow Replicas



Conclusion

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- Slowdowns hurt availability, need s-slowdown-tolerant RSMs
- Copilot: first 1-slowdown-tolerant protocol
 - Slowdown tolerance: proactive redundancy and fast takeovers
 - Optimizations: ping-pong batching and null dependency elimination
- Copilot's performance without slow replicas is competitive
- Copilot is the only protocol that tolerates any one slowdown